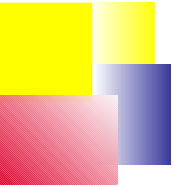


# Lightweight Formal Methods for the Development of High-Assurance Network Systems



**Assaf Kfoury**

with contributions from

**Azer Bestavros, Adam Bradley, Andrei Lapets, and Michael Ocean**

**iBench Initiative**

<http://www.cs.bu.edu/groups/ibench/>

**snBench**

<http://csr.bu.edu/snbench/>



**Computer Science**

# More Formal Methods ...

for the development of a rigorous discipline of *specification, analysis, programming and maintenance* of network systems

## 1. Compositional Analysis/Specification and its Benefits

(mostly with **Azer Bestavros**)

iBench Initiative – <http://www.cs.bu.edu/groups/ibench/>

## 2. An Application of Model Checking:

### Safe Composition of Arbitrary Network Protocols

(mostly with **Adam Bradley** and **Azer Bestavros**)

iBench Initiative – <http://www.cs.bu.edu/groups/ibench/>

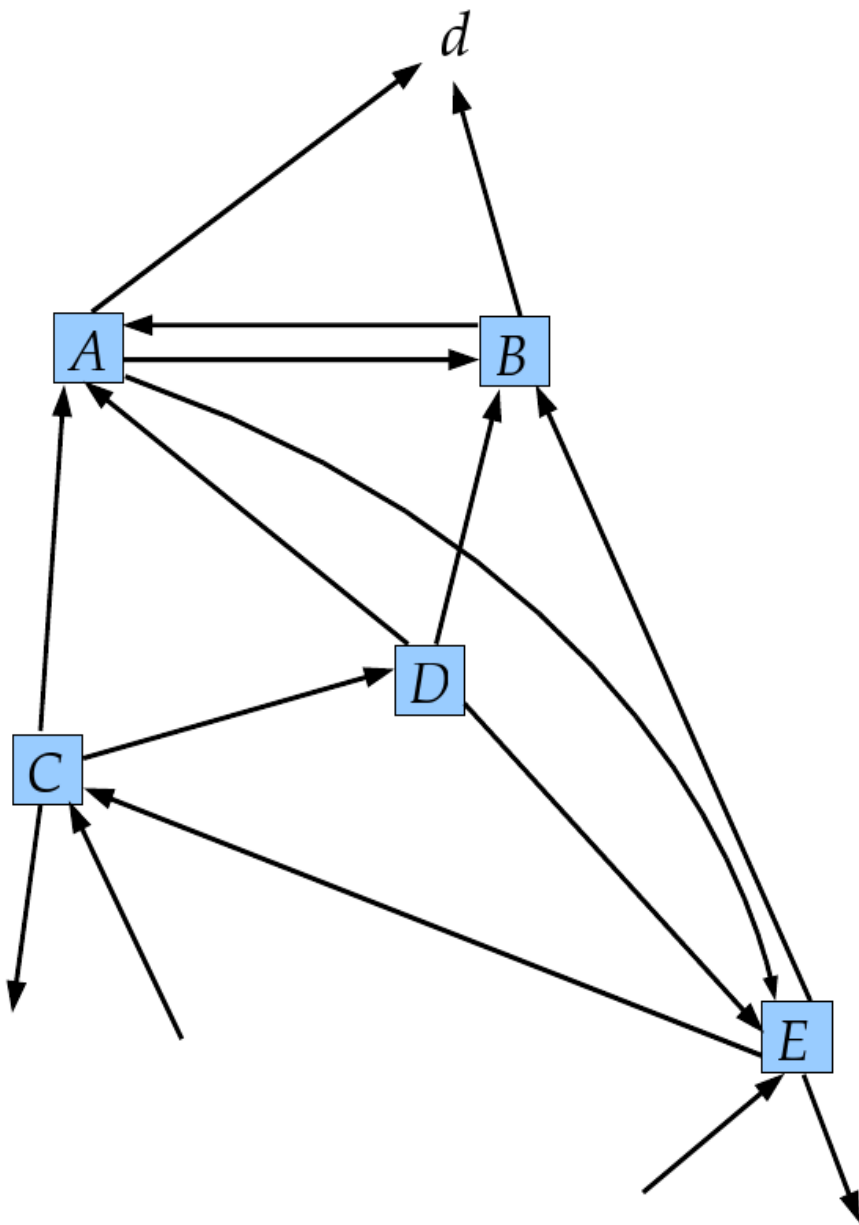
## 3. Resource Allocation in Sensor Networks Using a Strongly-Typed Domain-Specific Language

(mostly with **Michael Ocean** and **Azer Bestavros**)

snBench – <http://csr.bu.edu/snbench/>

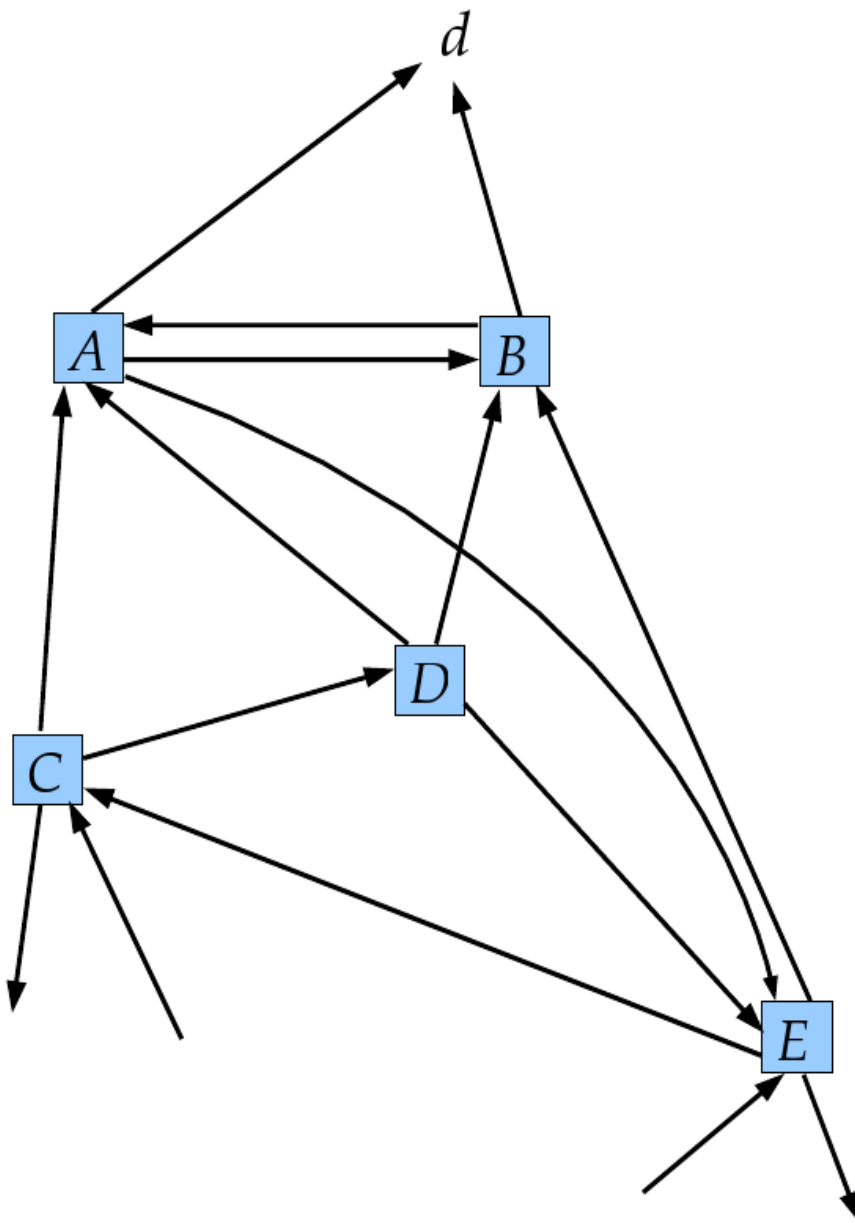
## 4. The Stable-Paths Problem and the Promise of an Automatic Lightweight Proof-Assistant

(mostly with **Kevin Donnelly** and **Andrei Lapets**)

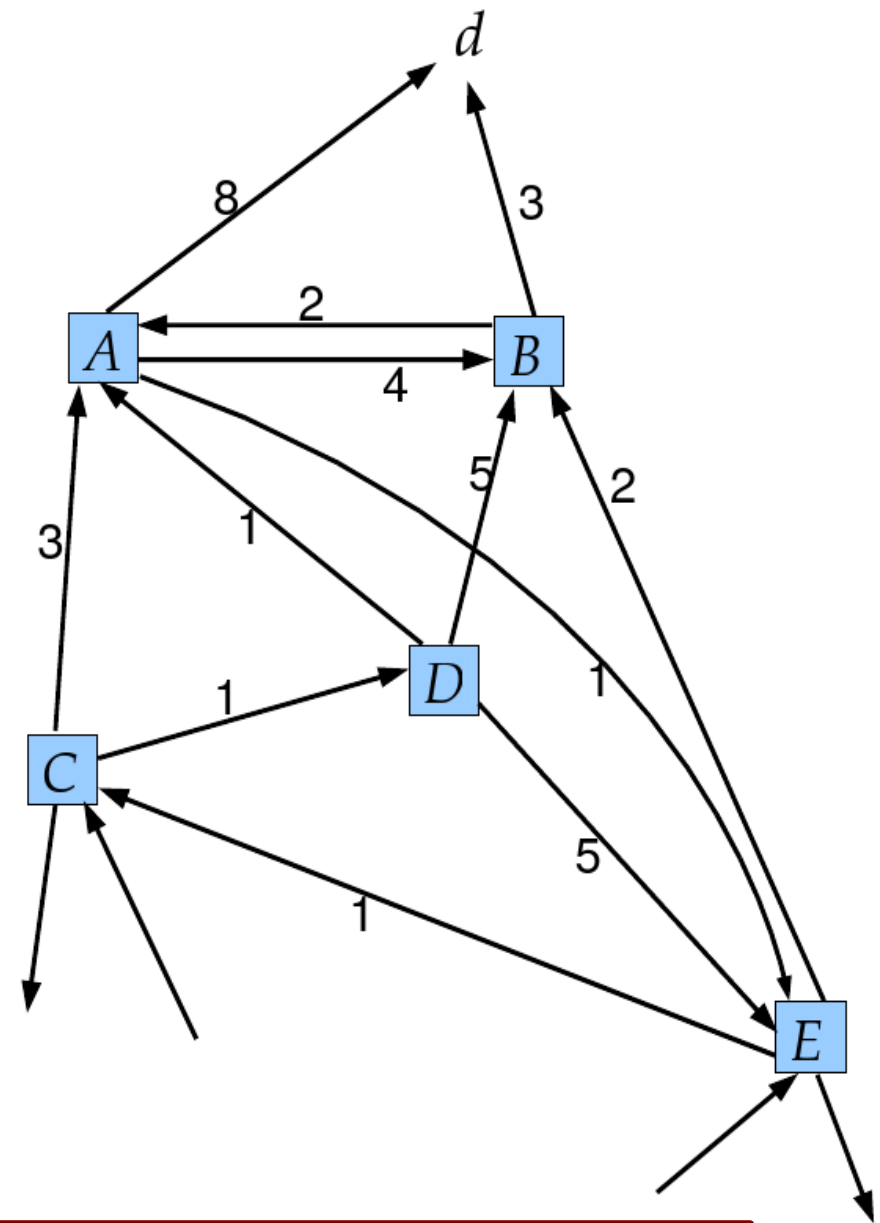


**network topology**

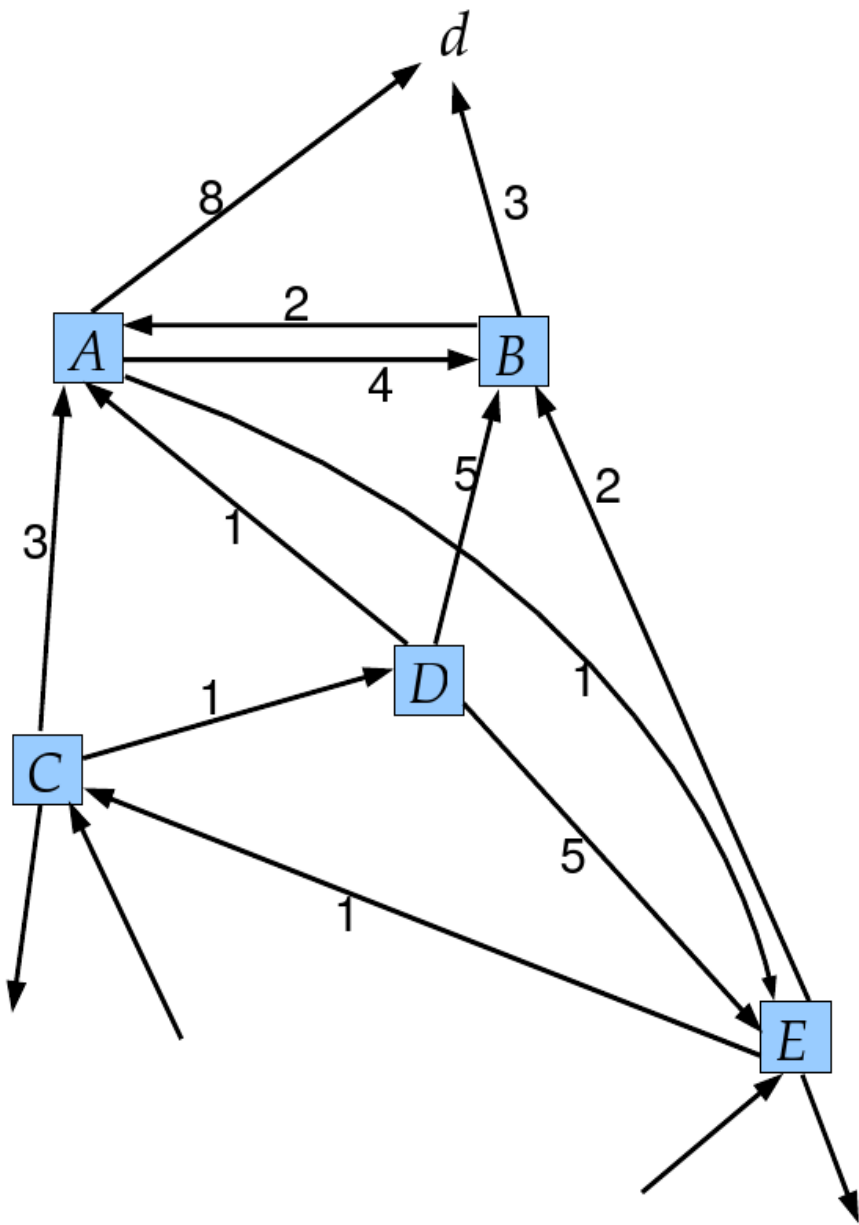
# one additive measure on links



network topology

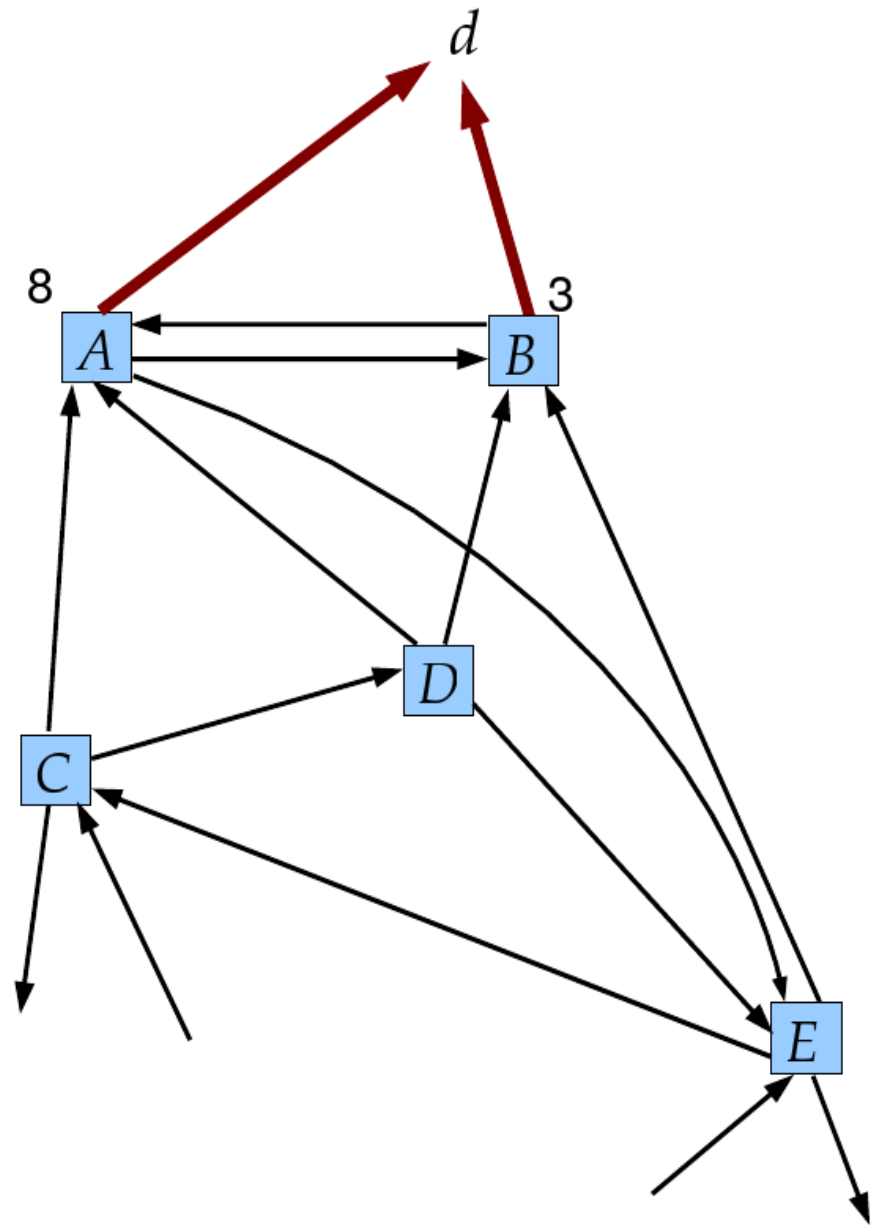


every node requests the  
"cheapest" path to  $d$   
every node communicates  
with its immediate peers only

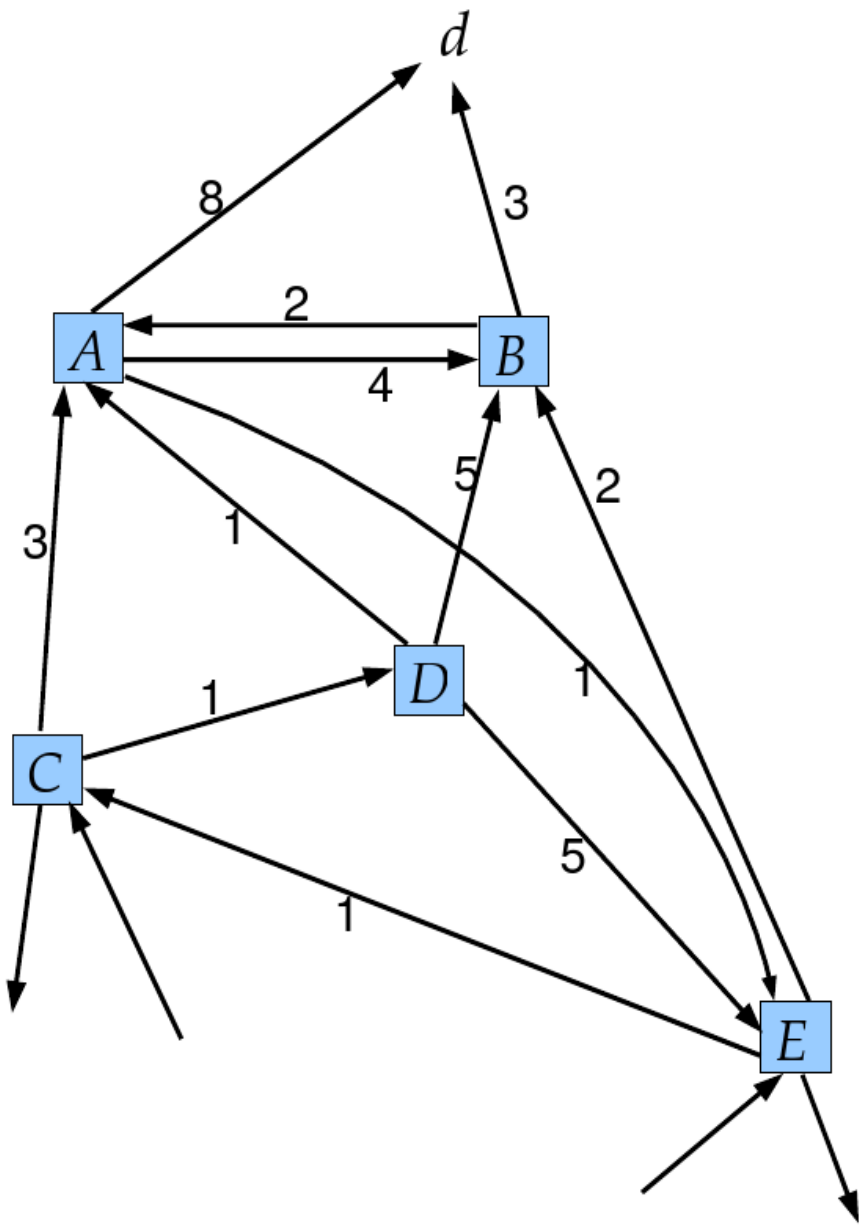


additive measure on links

at end of time = 1

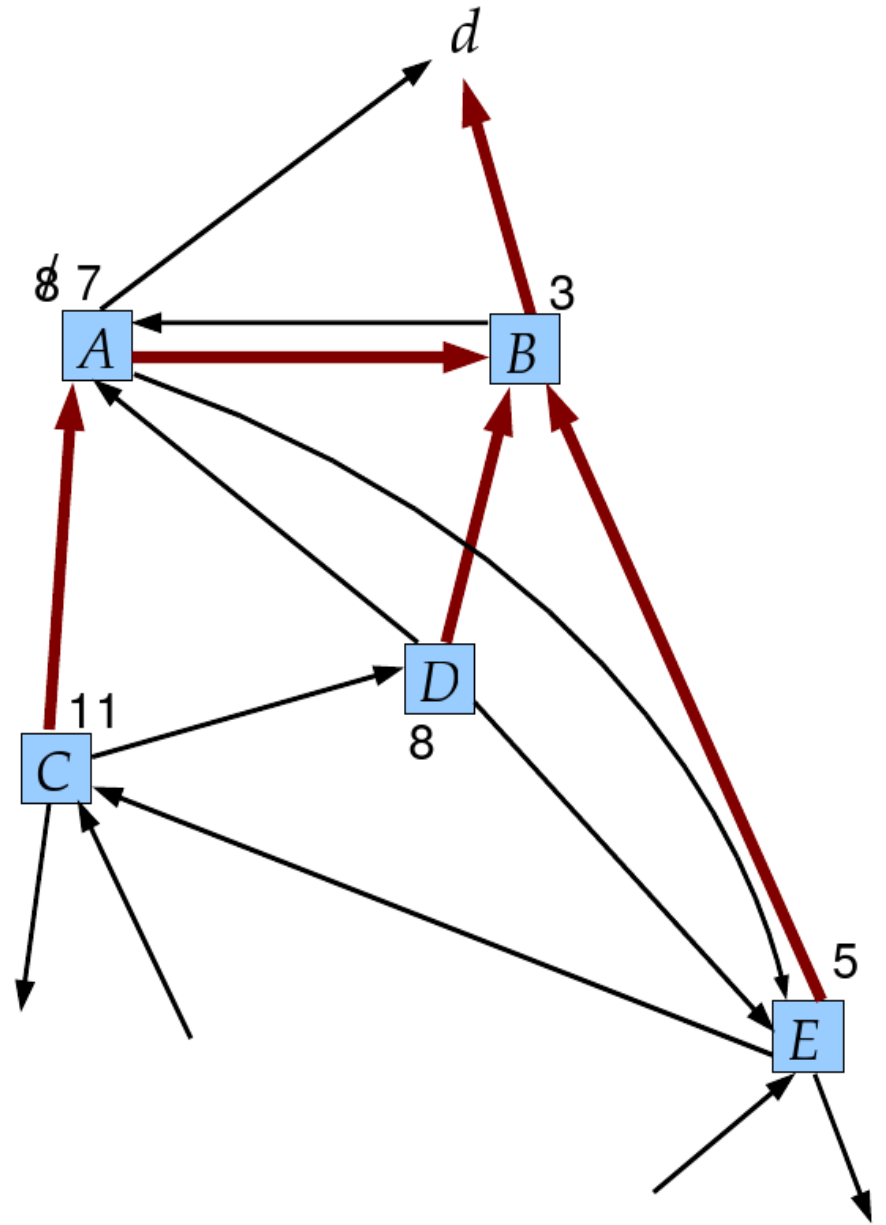


after  $d$  broadcasts to  $A$  and  $B$

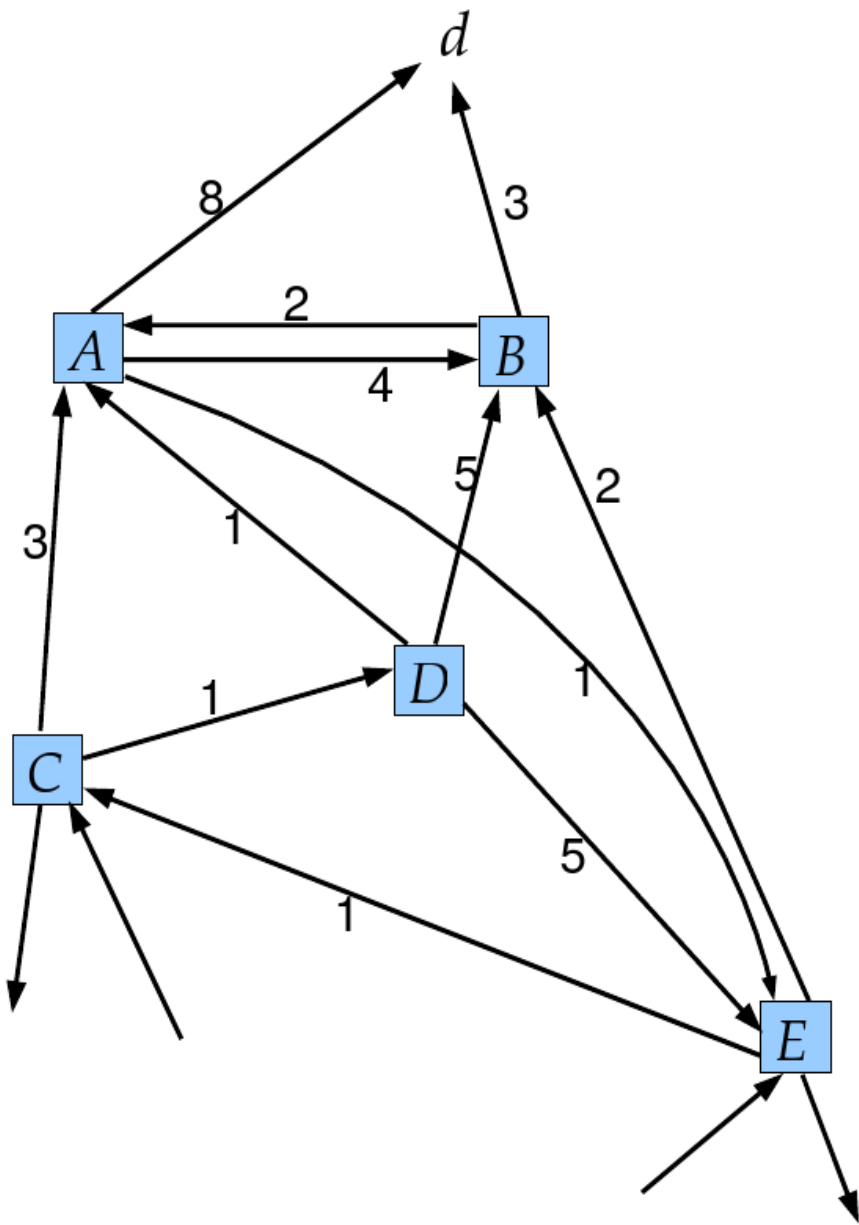


additive measure on links

at end of time = 2

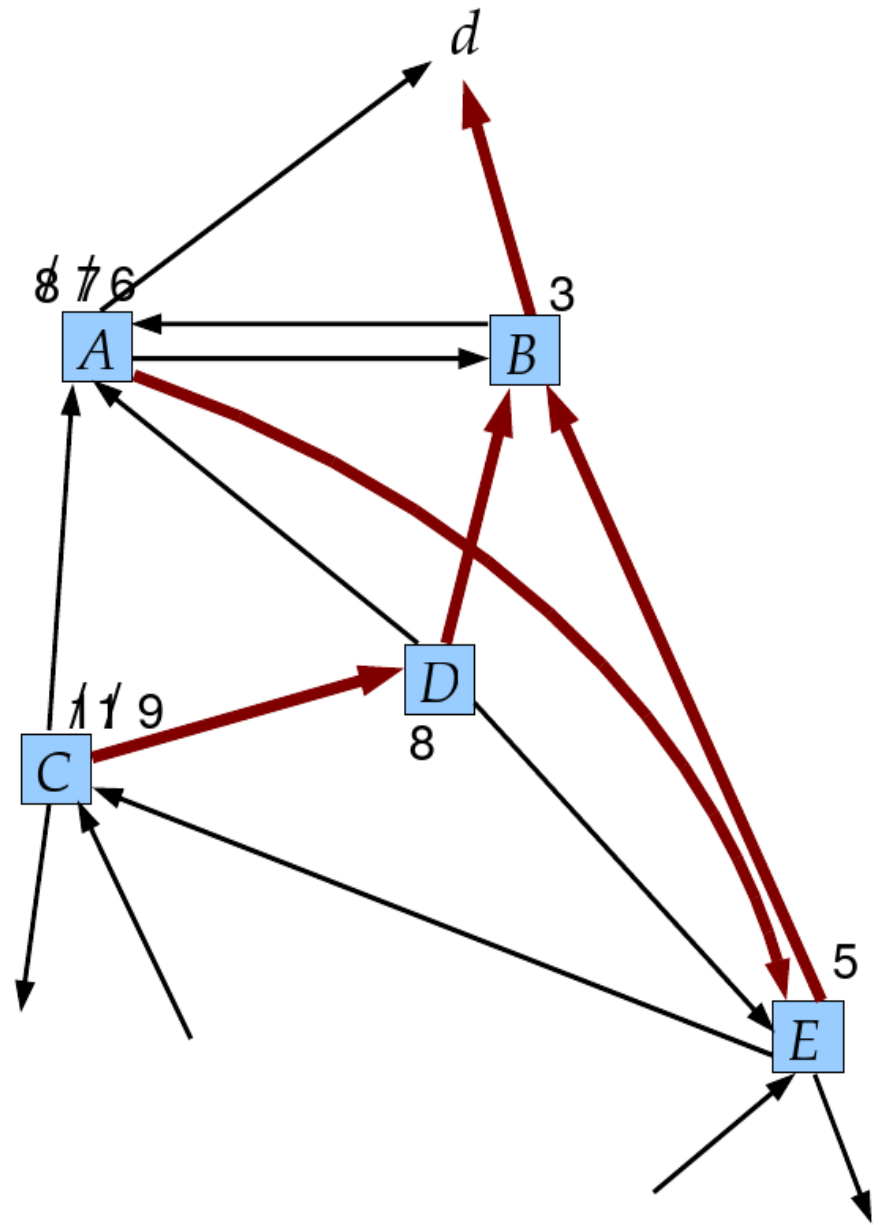


after  $A$  broadcasts to  $B, C, D$   
and  $B$  broadcasts to  $A, D, E$

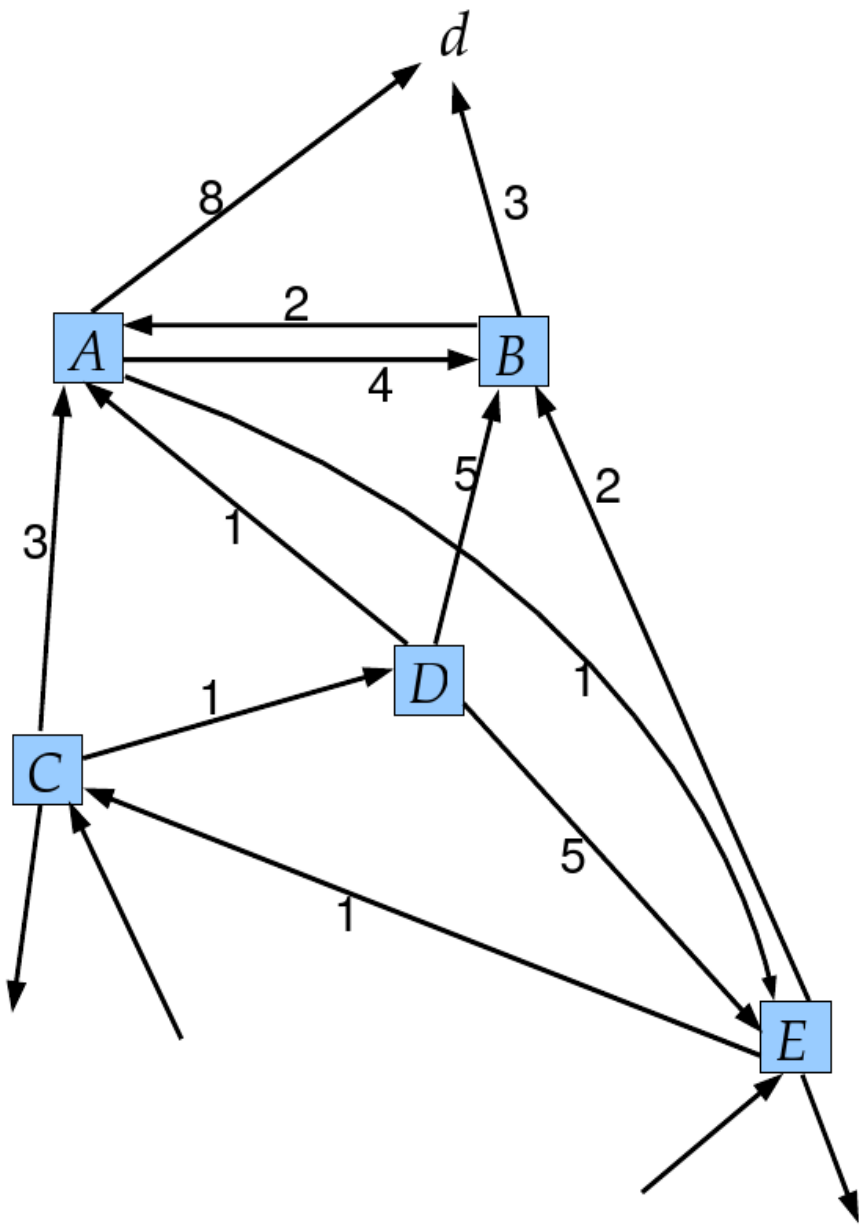


additive measure on links

at end of time = 3

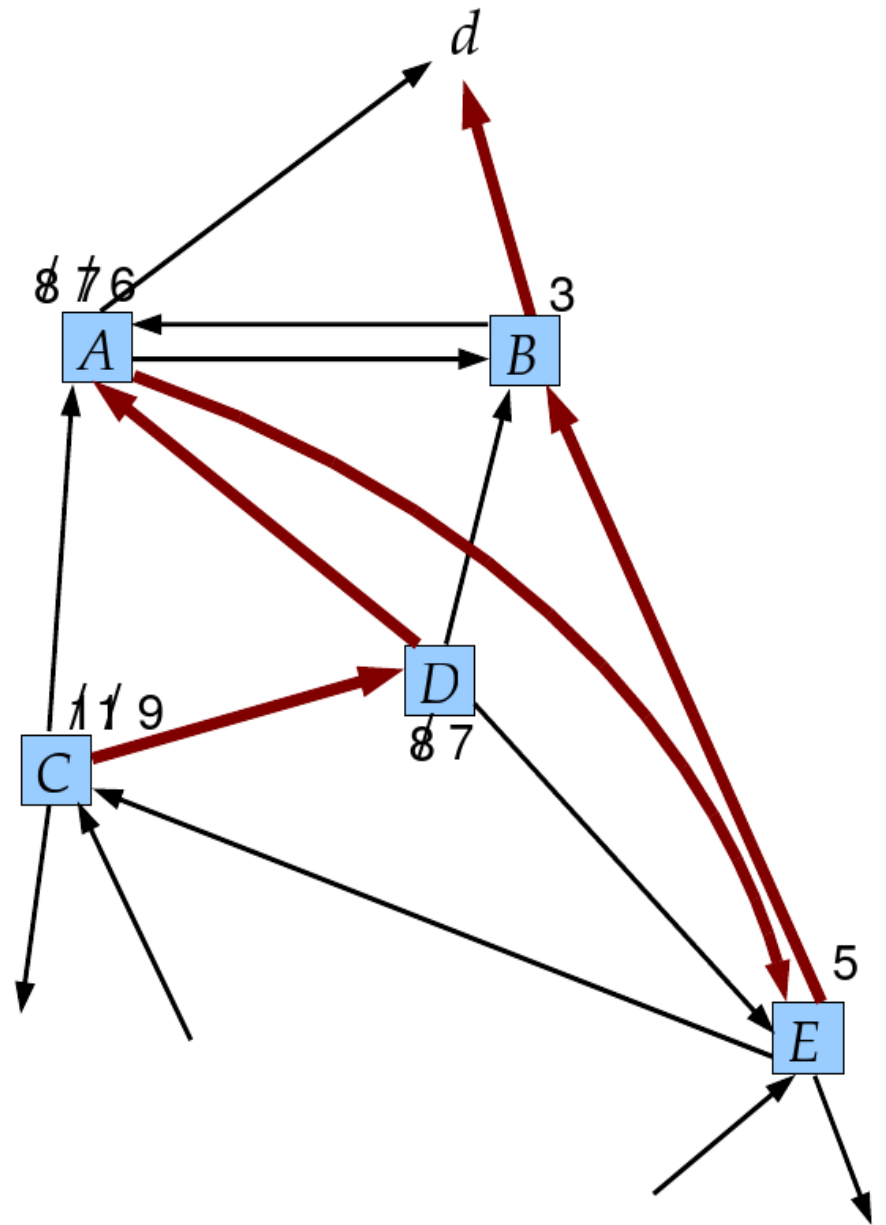


after  $C$  broadcasts to  $E, \dots$   
 $D$  broadcasts to  $C$   
 $E$  broadcasts to  $A, D, \dots$

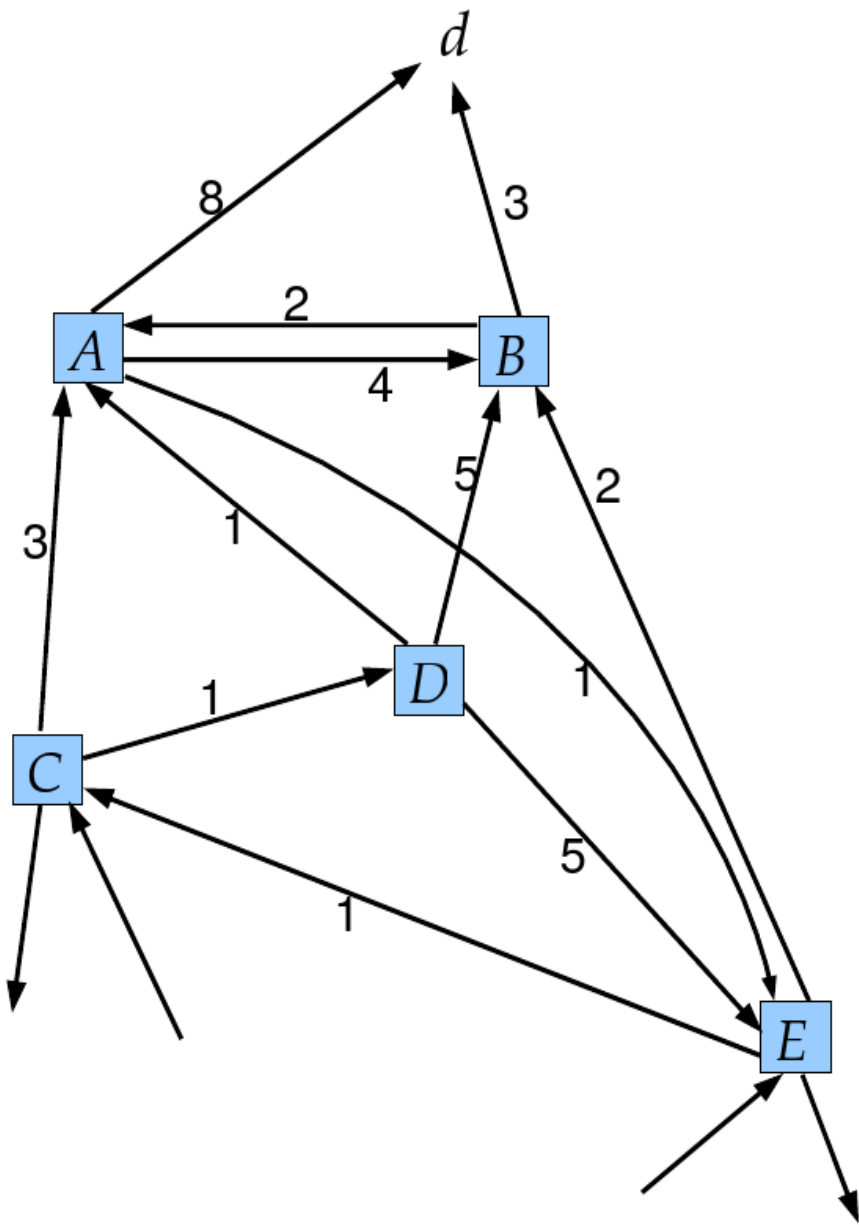


additive measure on links

at end of time = 4

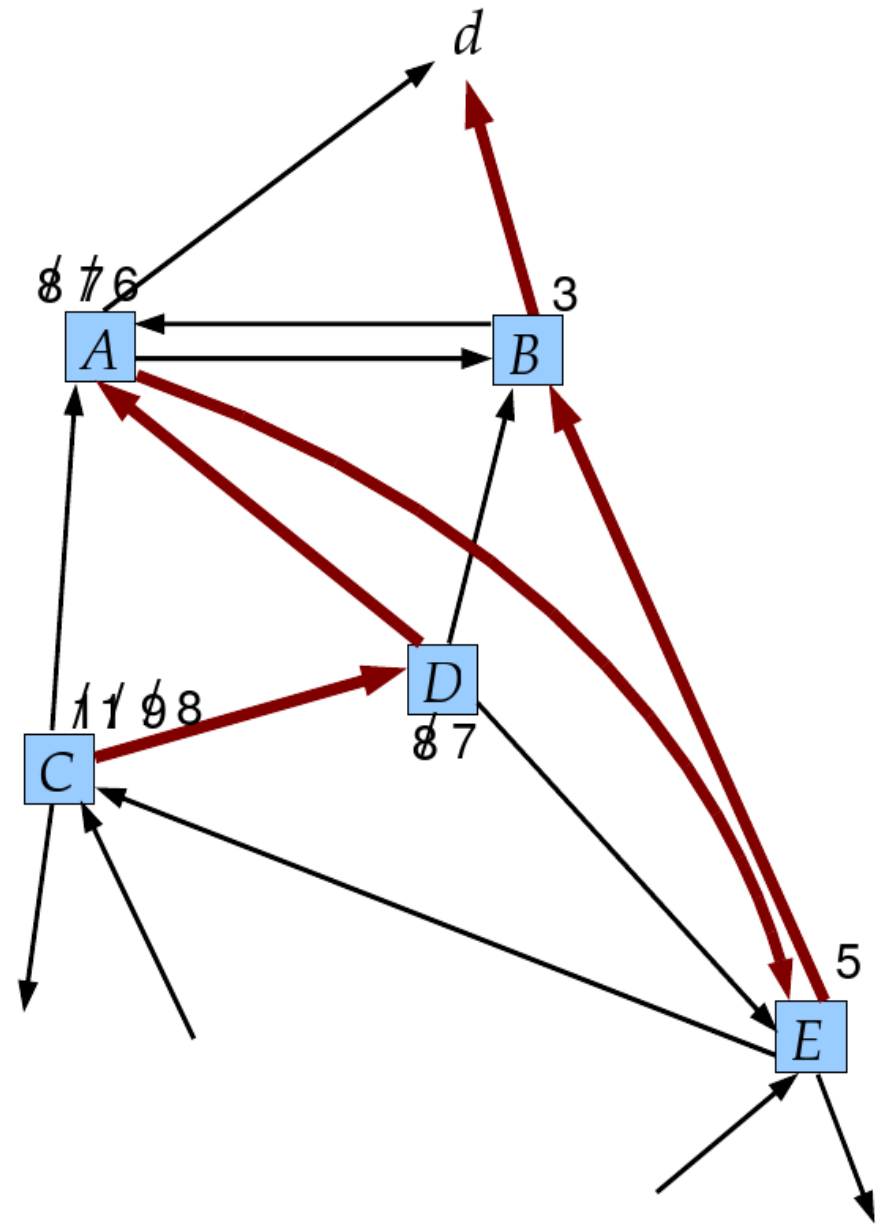


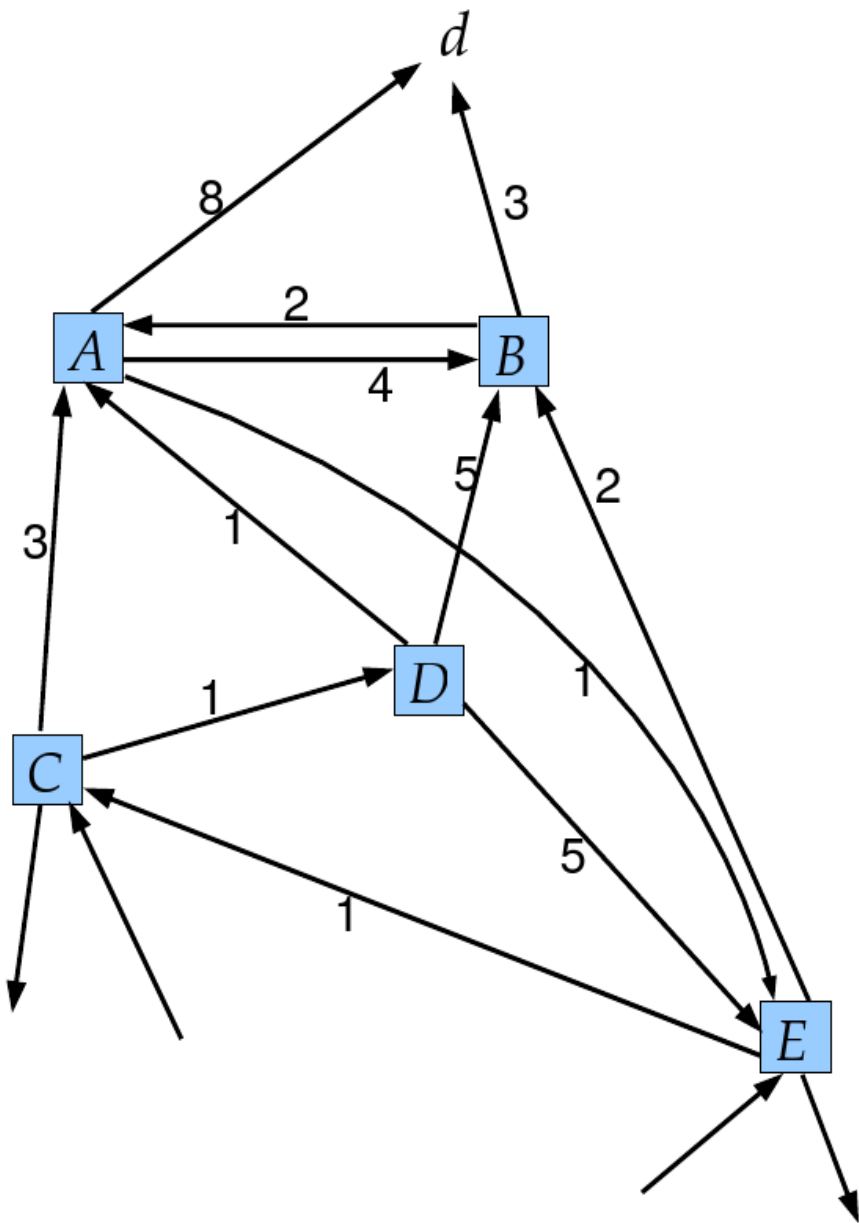




additive measure on links

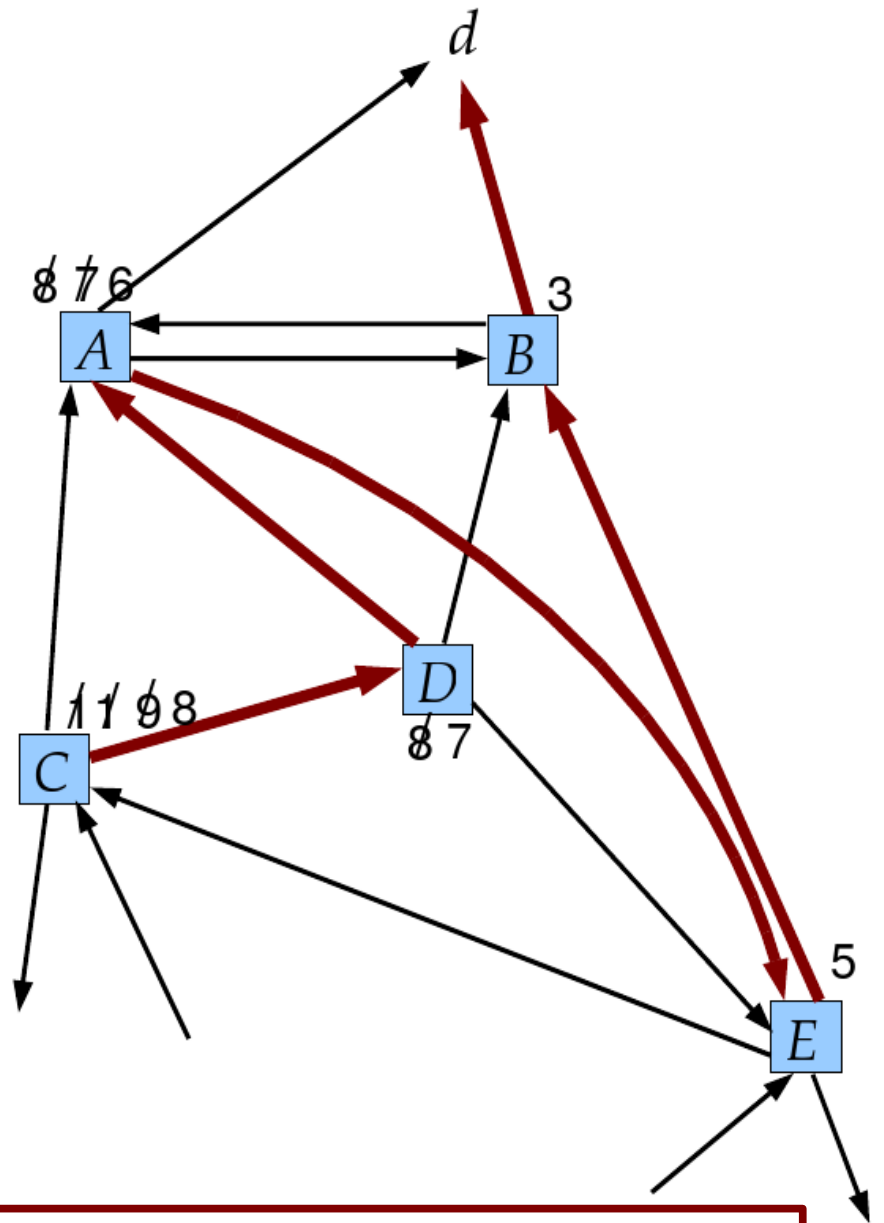
at end of time = 5





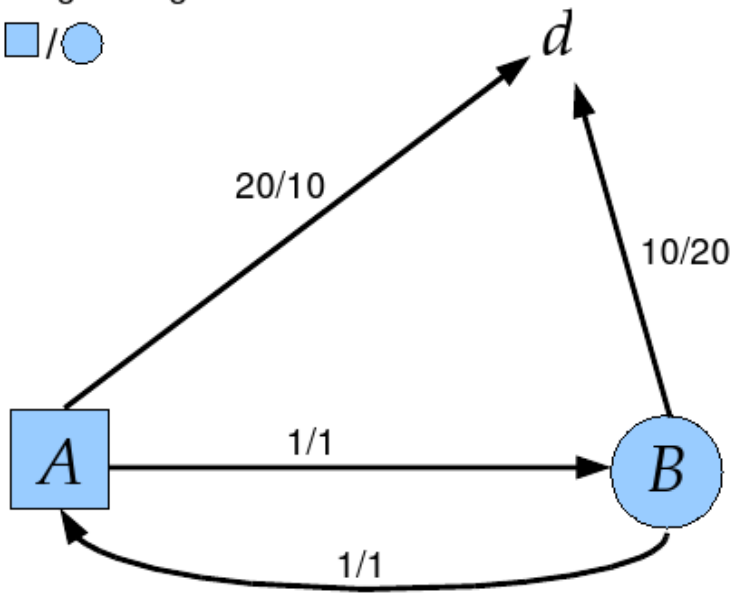
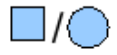
additive measure on links

at end of time = 5



**FACT.**  
 1. Each node finds a stable path to *d*.  
 2. If the network is finite, a MST rooted at *d* is constructed.

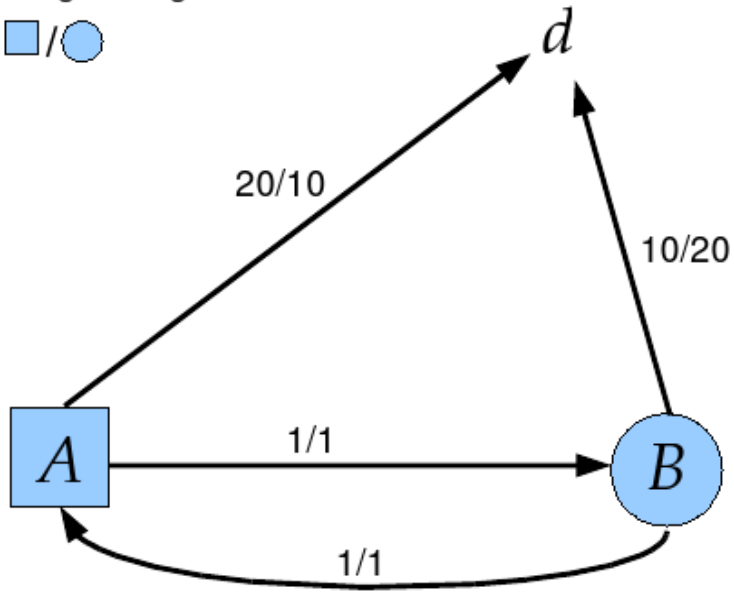
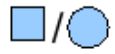
listing of weights:



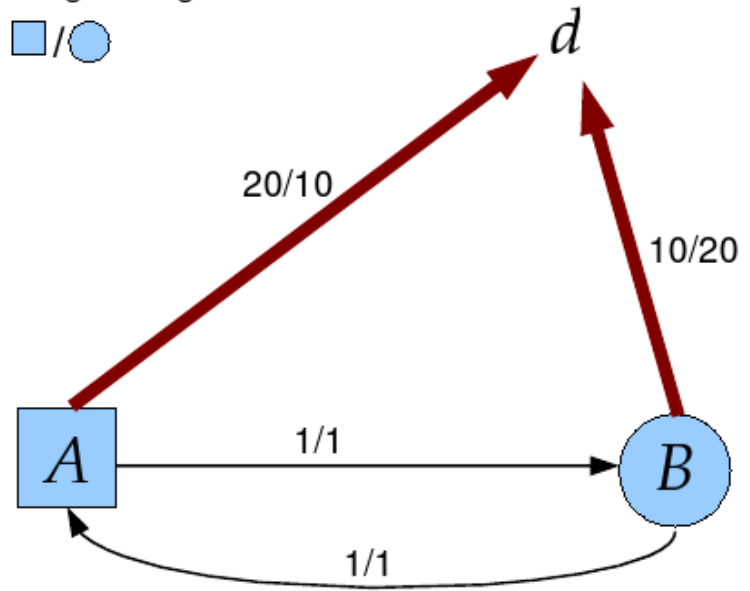
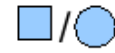
**two additive measures on links**

at end of time = 1

listing of weights:



listing of weights:

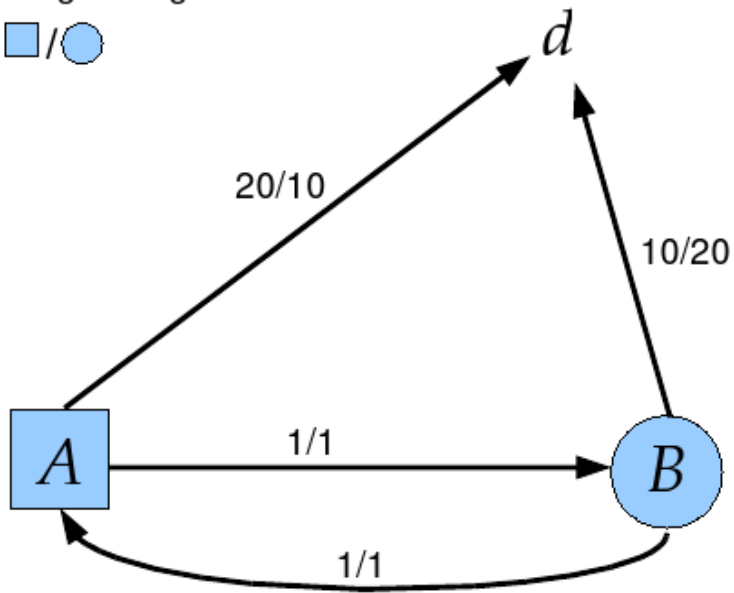


two additive measures on links

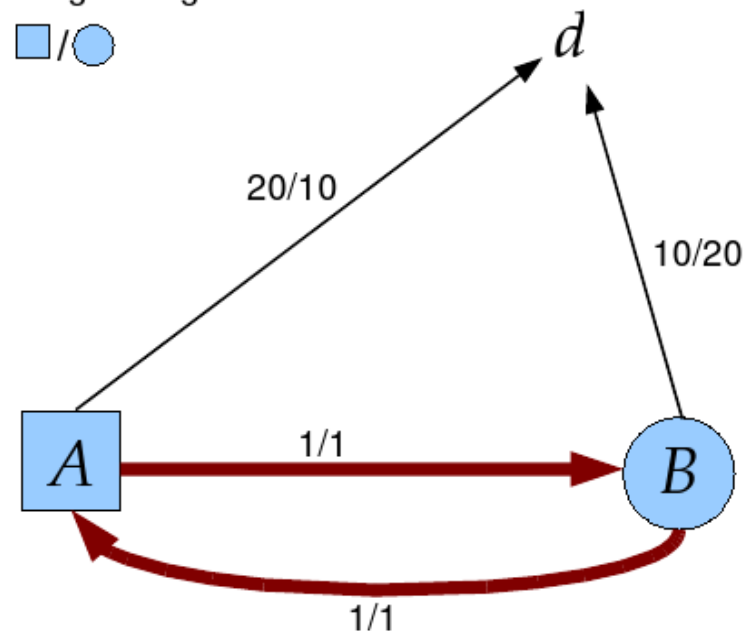
after *d* broadcasts to *A* and *B*

at end of time = 2

listing of weights:



listing of weights:

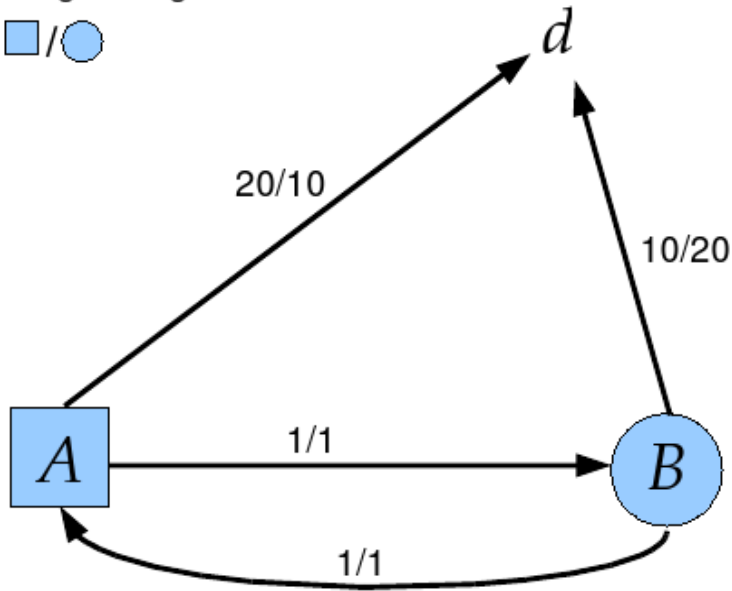


two additive measures on links

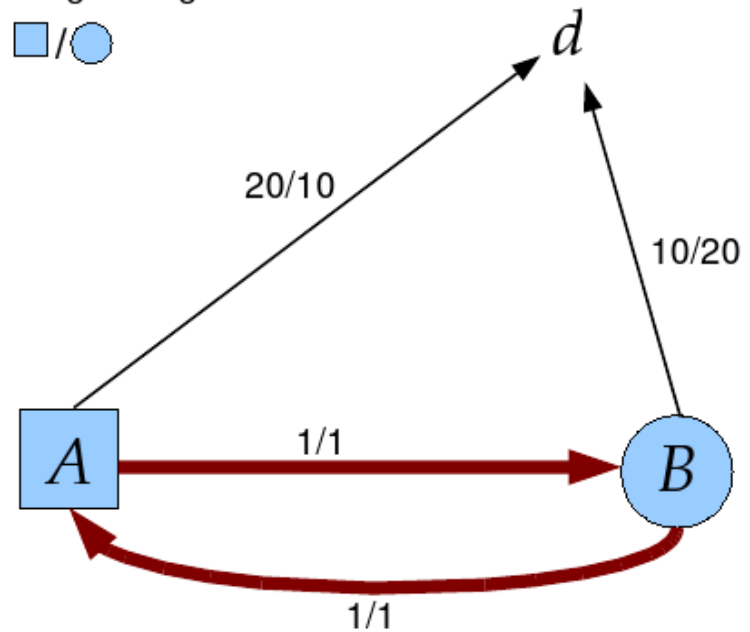
after **A** broadcasts to **B**  
and **B** broadcasts to **A**

at end of time = 2

listing of weights:



listing of weights:



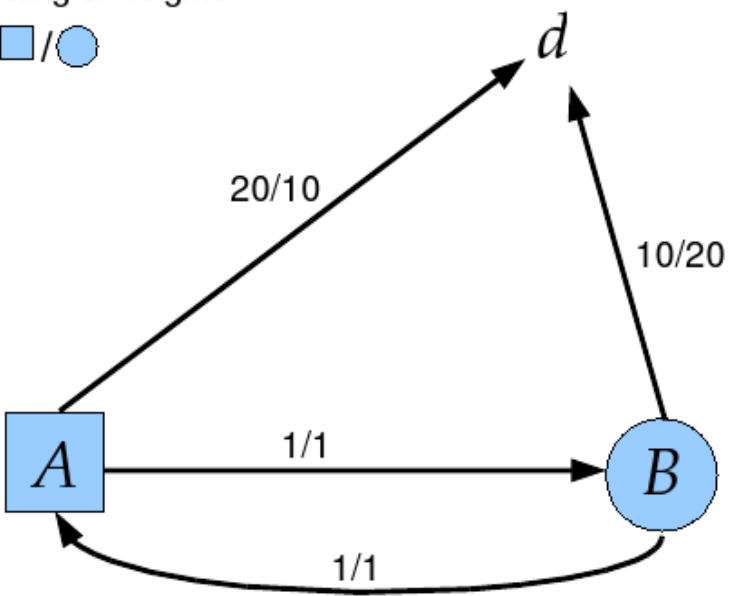
two additive measures on links

after **A** broadcasts to **B**  
and **B** broadcasts to **A**

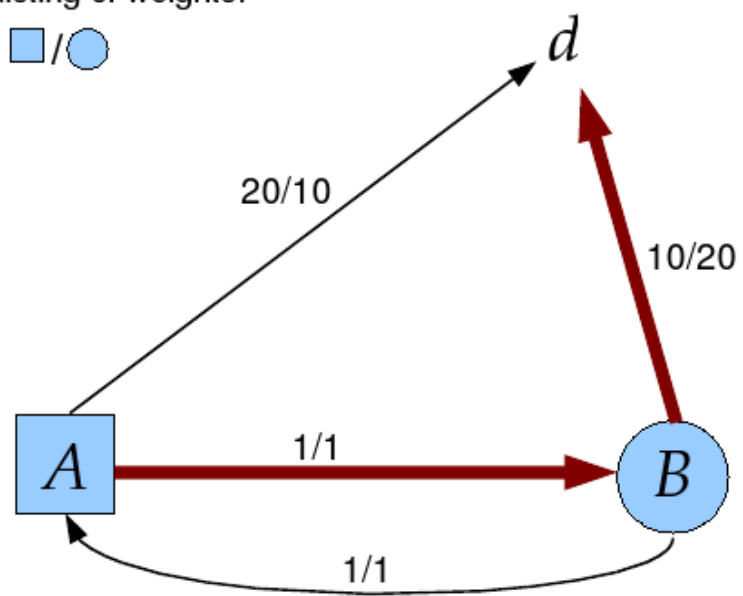
**FACT. No node (other than **d**)  
ever finds a stable path to **d**.**

**But there are stable configurations!**

listing of weights:



listing of weights:

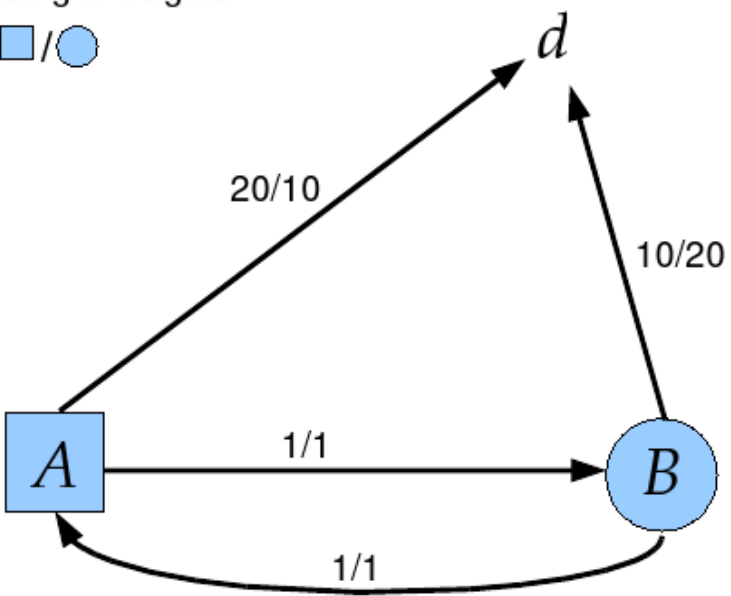


**stable configuration 1**

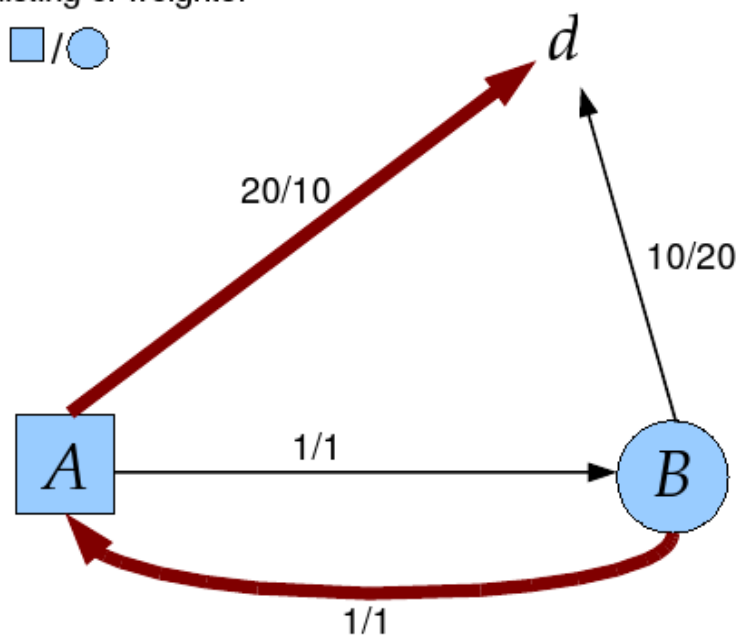
**two additive measures on links**

**But there are stable configurations!**

listing of weights:



listing of weights:



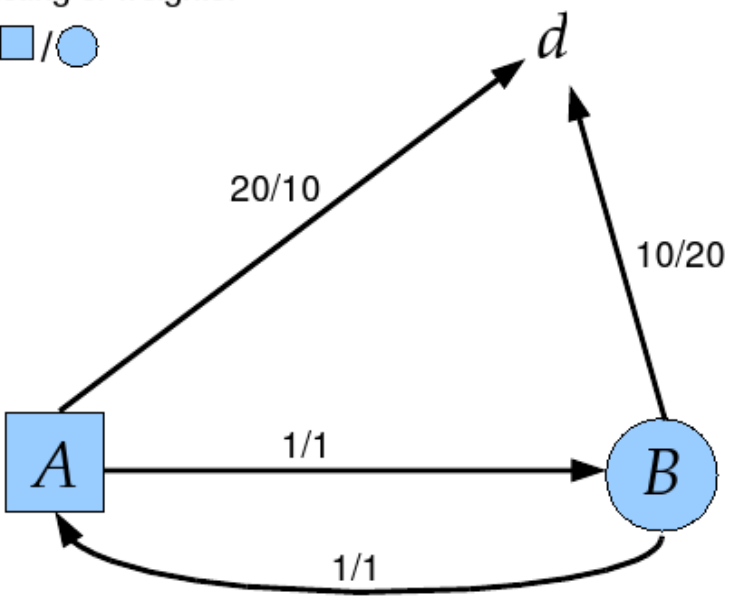
**stable configuration 2**

**two additive measures on links**

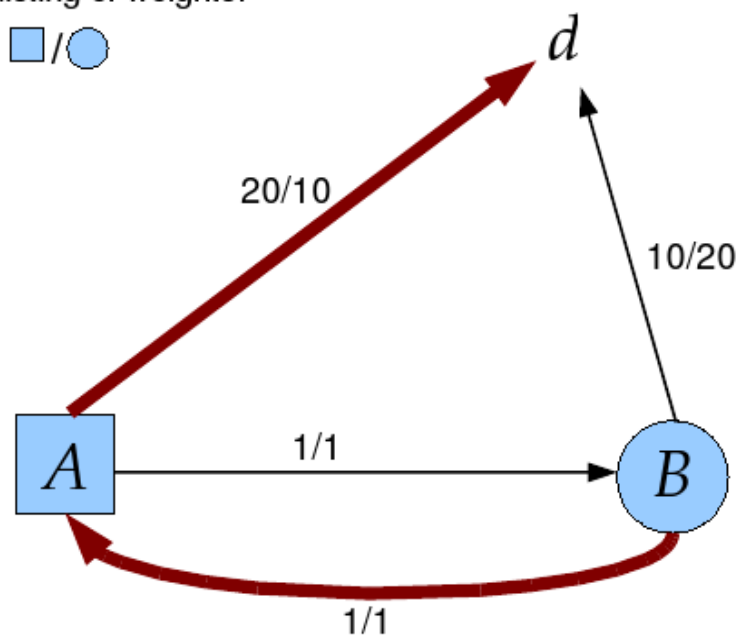


**But there are stable configurations!**

listing of weights:



listing of weights:



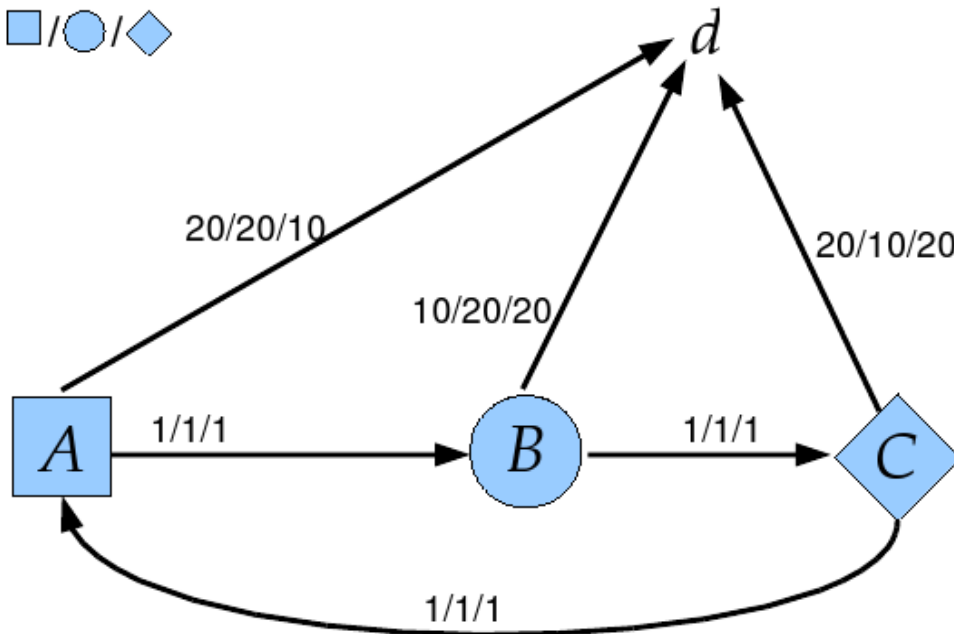
two additive measures on links

**stable configuration 2**

**Is synchrony the culprit?  
Perhaps there is an asynchronous procedure that will find stable paths ...**

**No!** There are networks where no node (other than *d*) can find a stable path to *d* – regardless of the method used.

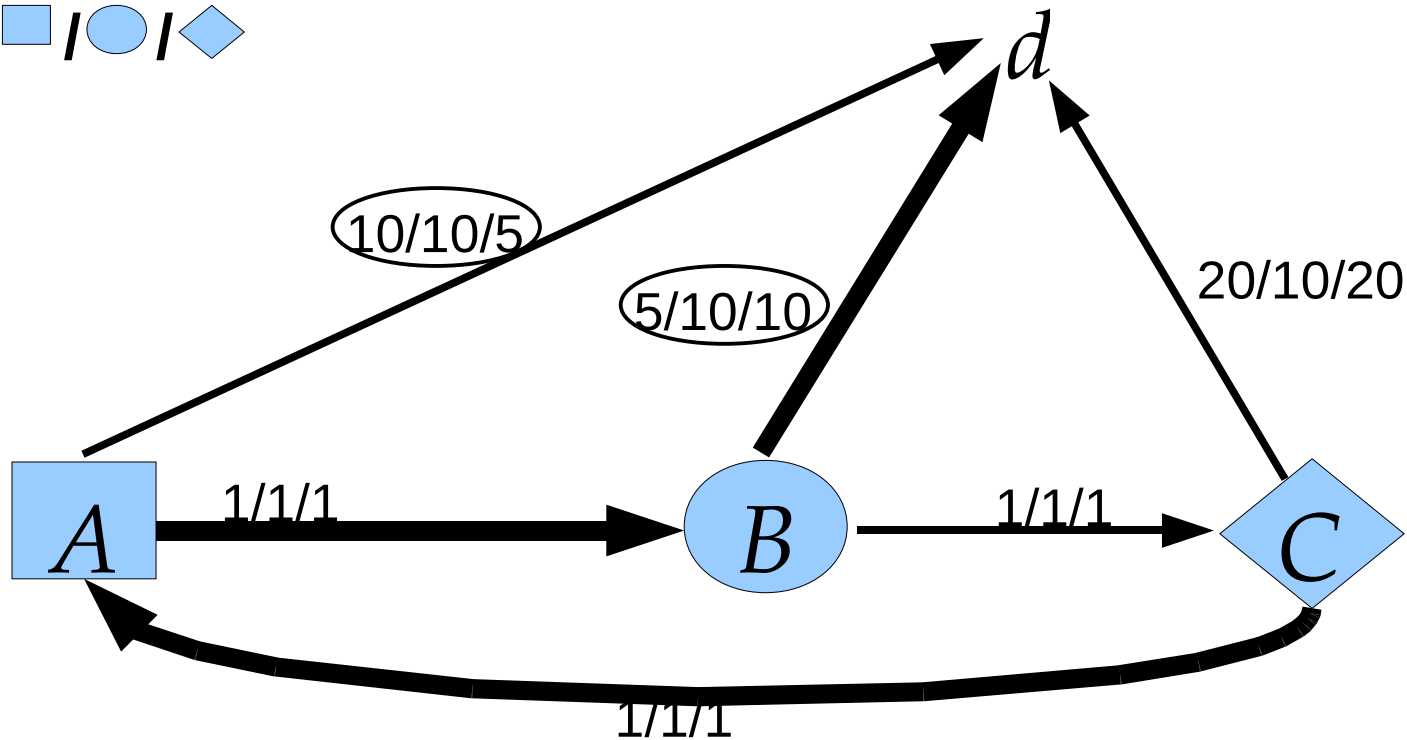
listing of weights:



- three additive measures on links
- inherently unstable network

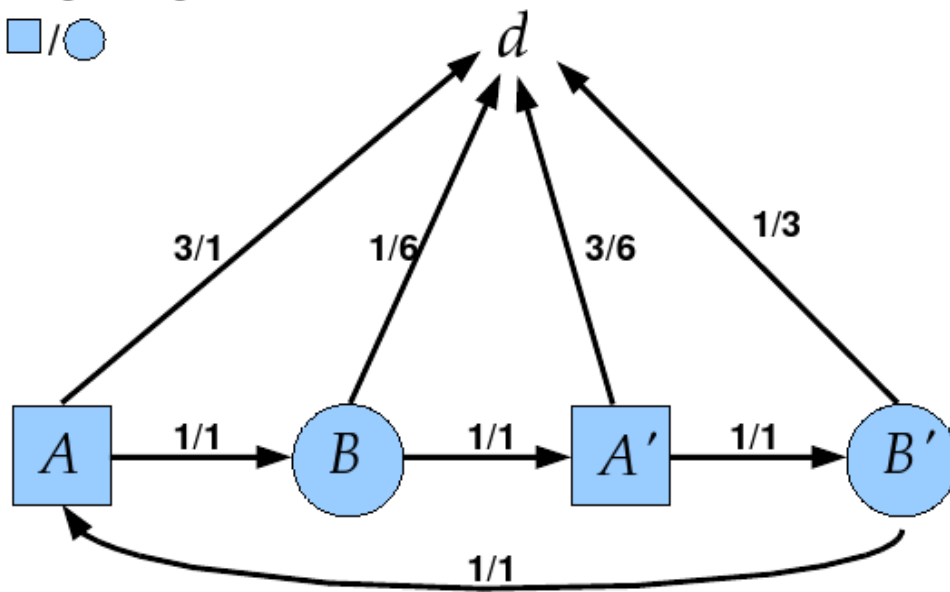
# Same network topology, slightly adjusted measures.

listing of weights:



- three additive measures on links
- exactly one stable configuration

listing of weights:



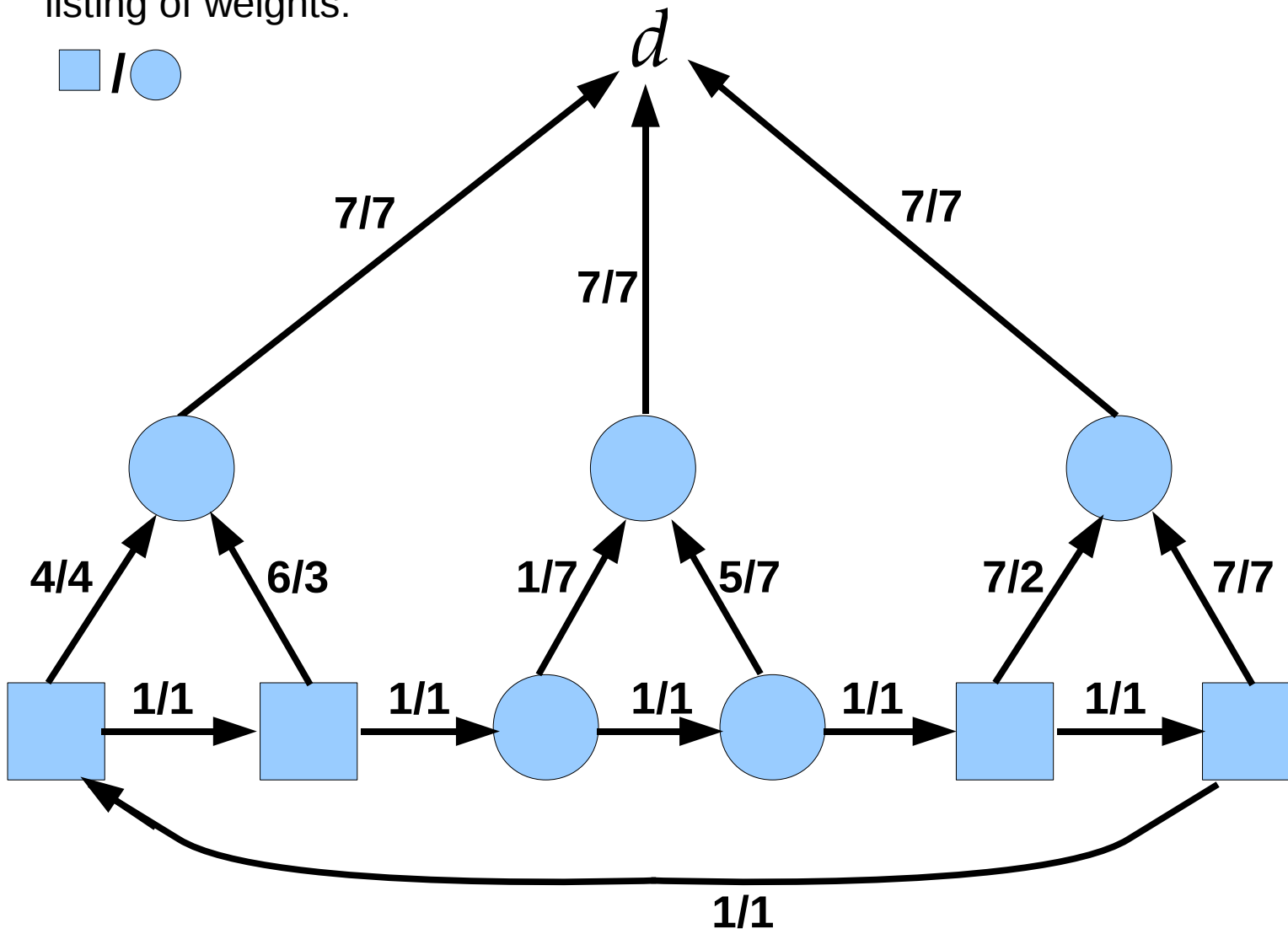
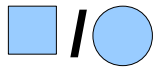
- **two additive measures on links**
- **inherently unstable network**

An instance of the **Stable Paths Problem (SPP)** is a network of **Autonomous Systems (AS)** each with a particular routing policies.

An instance of SPP is **finite** if the underlying network of AS's is finite.

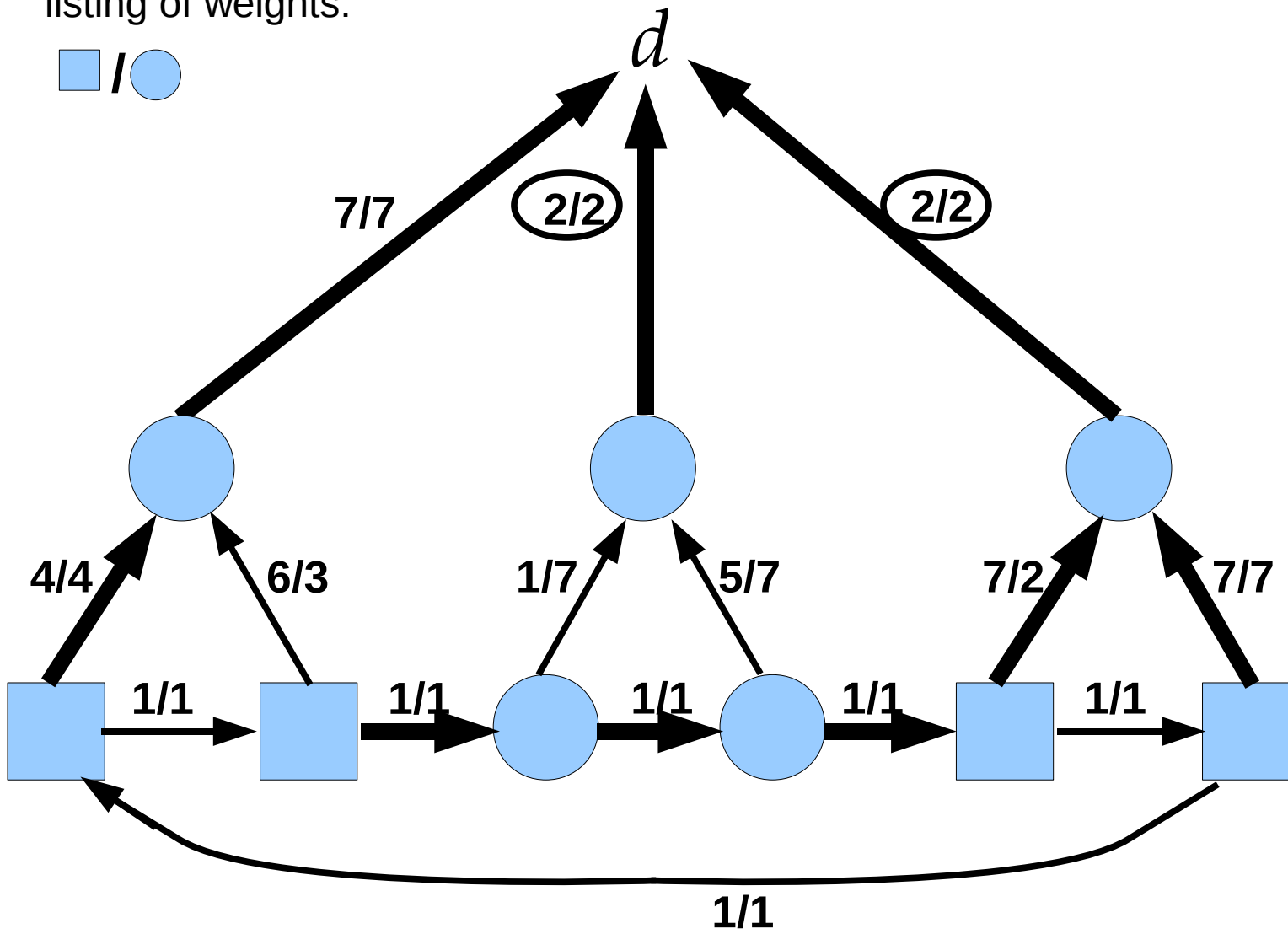
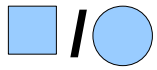
**THEOREM.** It is an NP-complete question whether a finite instance of SPP, with at least two routing policies, has a stable configuration.

listing of weights:



- two additive measures on links
- inherently unstable network

listing of weights:



- two additive measures on links
- exactly one stable configuration

## References

### ***1. On the Stable Paths Problem and a Restricted Variant***

**Kevin Donnelly and Assaf Kfoury**  
**BU Tech Report, 5 February 2008**

### **2. *ibis***

**lightweight logical framework and proof assistant**  
**Andrei Lapets**  
**<http://safre.org/>**



**Thank You!**

**Questions?**